

On the q -Analogue of Cauchy Matrices

Alessandro Neri

17-21 June 2019 - VUB



Universität
Zürich^{UZH}

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I am a friend of Finite Geometry!

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q -Analogs Model

Finite set	\rightsquigarrow	Finite dim vector space over \mathbb{F}_q
Element	\rightsquigarrow	1-dim subspace
\emptyset	\rightsquigarrow	$\{0\}$
Cardinality	\rightsquigarrow	Dimension
Intersection	\rightsquigarrow	Intersection
Union	\rightsquigarrow	Sum

Examples

1. Binomials and q -binomials:

$$\binom{n}{k} = \prod_{i=0}^{k-1} \frac{n-i}{k-i}, \quad \binom{n}{k}_q = \prod_{i=0}^{k-1} \frac{1-q^{n-i}}{1-q^{k-i}}.$$

2. (Chu)-Vandermonde and q -Vandermonde identity:

$$\binom{m+n}{k} = \sum_{j=0}^k \binom{m}{j} \binom{n}{k-j}, \quad \binom{m+n}{k}_q = \sum_{j=0}^k \binom{m}{k-j}_q \binom{n}{j}_q q^{j(m-k+j)}.$$

3. Polynomials and q -polynomials:

$$a_0 + a_1 x + \dots + a_k x^k, \quad a_0 x + a_1 x^q + \dots + a_k x^{q^k}.$$

4. Gamma and q -Gamma functions:

$$\Gamma(z+1) = z\Gamma(z), \quad \Gamma_q(x+1) = \frac{1-q^x}{1-q} \Gamma_q(x).$$

Vandermonde Matrix

Let $k \leq n$.

$$V = \begin{pmatrix} 1 & 1 & \dots & 1 \\ \alpha_1 & \alpha_2 & \dots & \alpha_n \\ \alpha_1^2 & \alpha_2^2 & \dots & \alpha_n^2 \\ \vdots & \vdots & & \vdots \\ \alpha_1^{k-1} & \alpha_2^{k-1} & \dots & \alpha_n^{k-1} \end{pmatrix}$$

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For $n = k$, $\det(V) \neq 0$ if and only if the α_i 's are all distinct.

$$|\{\alpha_1, \dots, \alpha_n\}| = n.$$

In particular, all the $k \times k$ minors of V are non-zero.

Moore Matrix

Let $k \leq n$.

$$G = \begin{pmatrix} g_1 & g_2 & \cdots & g_n \\ g_1^q & g_2^q & \cdots & g_n^q \\ \vdots & \vdots & & \vdots \\ g_1^{q^{k-1}} & g_2^{q^{k-1}} & \cdots & g_n^{q^{k-1}} \end{pmatrix},$$

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For $n = k$, $\det(G) \neq 0$ if and only if the g_i 's are all \mathbb{F}_q -linearly independent.

$$\dim_{\mathbb{F}_q} \langle g_1, \dots, g_n \rangle_{\mathbb{F}_q} = n.$$

In particular, for every $M \in \mathrm{GL}_n(\mathbb{F}_q)$, all the $k \times k$ minors of GM are non-zero.

Moore Matrix

Let $k \leq n$,

$$\sigma : x \mapsto x^q$$

$$G = \begin{pmatrix} g_1 & g_2 & \dots & g_n \\ \sigma(g_1) & \sigma(g_2) & \dots & \sigma(g_n) \\ \vdots & \vdots & & \vdots \\ \sigma^{k-1}(g_1) & \sigma^{k-1}(g_2) & \dots & \sigma^{k-1}(g_n) \end{pmatrix},$$

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Moore matrix

1. Dickson used the Moore matrix for finding the modular invariants of the general linear group over a finite field.
2. It is widely used in the study of normal bases of finite fields.
- 3.

$$\det(G) = \prod_{1 \leq i \leq n} \prod_{c_1, \dots, c_{i-1} \in \mathbb{F}_q} (c_1 g_1 + \dots + c_{i-1} g_{i-1} + g_i).$$

Generalized Cauchy Matrix

- (1) $x_1, \dots, x_r \in \mathbb{F}_q$ pairwise distinct,
- (2) $y_1, \dots, y_s \in \mathbb{F}_q$ pairwise distinct,
- (3) $y_1, \dots, y_s \in \mathbb{F}_q \setminus \{x_1, \dots, x_r\}$,
- (4) $c_1, \dots, c_r, d_1, \dots, d_s \in \mathbb{F}_q^*$.

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The matrix $X \in \mathbb{F}_q^{r \times s}$ defined by

$$X_{i,j} = \frac{c_i d_j}{x_i - y_j}$$

is called Generalized Cauchy (GC) Matrix.

For $j \leq \min\{r, s\}$, all the $j \times j$ minors are non-zero.

Question

How to define a q -analogue
of Cauchy matrices?

Generalized Reed-Solomon Codes

[Reed, Solomon '60]

- $\mathbb{F}_q[x]_{<k} := \{f(x) \in \mathbb{F}_q[x] \mid \deg f < k\}$.

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- $b_1, \dots, b_n \in \mathbb{F}_q^*$
- $\mathcal{C} = \{(b_1 f(\alpha_1), b_2 f(\alpha_2), \dots, b_n f(\alpha_n)) \mid f \in \mathbb{F}_q[x]_{<k}\}$

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Canonical Generator Matrix for GRS codes

Consider the canonical monomial basis $\{1, x, \dots, x^{k-1}\}$ and evaluate it.
We get the following generator matrix

Weighted Vandermonde (WV) matrix

$$\begin{pmatrix} 1 & 1 & \dots & 1 \\ \alpha_1 & \alpha_2 & \dots & \alpha_n \\ \alpha_1^2 & \alpha_2^2 & \dots & \alpha_n^2 \\ \vdots & \vdots & & \vdots \\ \alpha_1^{k-1} & \alpha_2^{k-1} & \dots & \alpha_n^{k-1} \end{pmatrix} \text{diag}(\mathbf{b})$$

GRS Codes and Cauchy Matrices

1. Every $[n, k]_q$ code has many generator matrices $G \in \mathbb{F}_q^{k \times n}$.
2. Every $[n, k]_q$ code has a unique generator matrix in Reduced Row Echelon Form (RREF).

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3. If \mathcal{C} is MDS the generator matrix in RREF is of the form $(I_k \mid X)$.
4. GRS codes are MDS.

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Theorem [Roth, Seroussi '85]

There is a 1-1 correspondence between GC matrices and GRS codes given by

$$X \longleftrightarrow \text{rowsp}(I_k \mid X).$$

General Setting

- ⌚ \mathbb{F}_{q^m} extension field of degree m of a finite field \mathbb{F}_q .
- ⌚ $\mathbb{F}_{q^m} \cong \mathbb{F}_q^m$ as vector spaces over \mathbb{F}_q .
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Rank Distance

The rank distance d_R on $\mathbb{F}_q^{m \times n}$ is defined by

$$d_R(X, Y) := \text{rk}(X - Y), \quad X, Y \in \mathbb{F}_q^{m \times n}.$$

The rank distance d_R on $\mathbb{F}_{q^m}^n$ is defined by

$$d_R(u, v) := \dim \langle u_1 - v_1, u_2 - v_2, \dots, u_n - v_n \rangle_{\mathbb{F}_q}.$$

A rank metric code \mathcal{C} is a subset of $\mathbb{F}_{q^m}^n$ equipped with the rank distance.

Linearized Polynomials and Gabidulin Codes

[Delsarte '78], [Gabidulin '85], [Roth '91], [Kshevetskiy, Gabidulin '05]

- $\sum_{i=0}^{m-1} f_i x^{q^i}$ a linearized polynomial over \mathbb{F}_{q^m} ,
- $\mathcal{G}_k := \left\{ f_0 x + f_1 x^q + \dots + f_{k-1} x^{q^{k-1}} \mid f_i \in \mathbb{F}_{q^m} \right\}.$

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- $g_1, \dots, g_n \in \mathbb{F}_{q^m}$ linearly independent over \mathbb{F}_q
 - $\mathcal{C} = \{(f(g_1), f(g_2), \dots, f(g_n)) \mid f \in \mathcal{G}_k\}$

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\mathcal{C} is the Gabidulin code $\mathcal{G}_k(g_1, \dots, g_n)$

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Canonical Generator Matrix for $\mathcal{G}_k(g_1, \dots, g_n)$

Let

$$\begin{aligned}\sigma : \mathbb{F}_{q^m} &\longrightarrow \mathbb{F}_{q^m} \\ z &\longmapsto z^q\end{aligned}$$

be the q -Frobenius automorphism.

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s-Moore Matrix

$$\left(\begin{array}{cccc} g_1 & g_2 & \dots & g_n \\ \sigma(g_1) & \sigma(g_2) & \dots & \sigma(g_n) \\ \vdots & \vdots & & \vdots \\ \sigma^{k-1}(g_1) & \sigma^{k-1}(g_2) & \dots & \sigma^{k-1}(g_n) \end{array} \right),$$

Hamming vs Rank Distance: Recap

d_H
 d_R

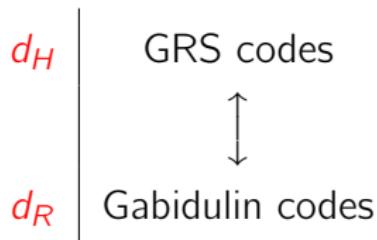
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d_H

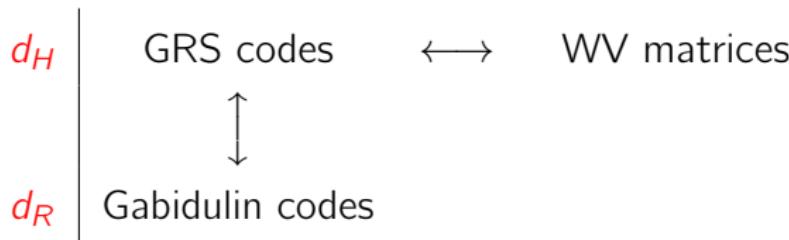
GRS codes

d_R

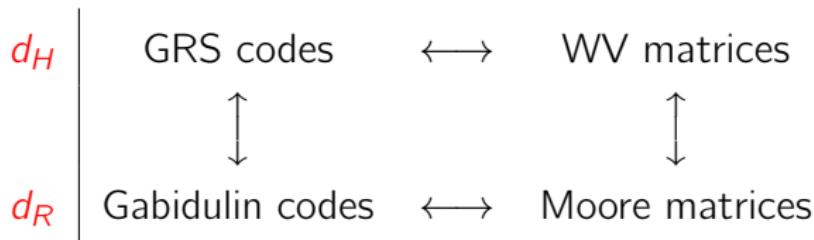
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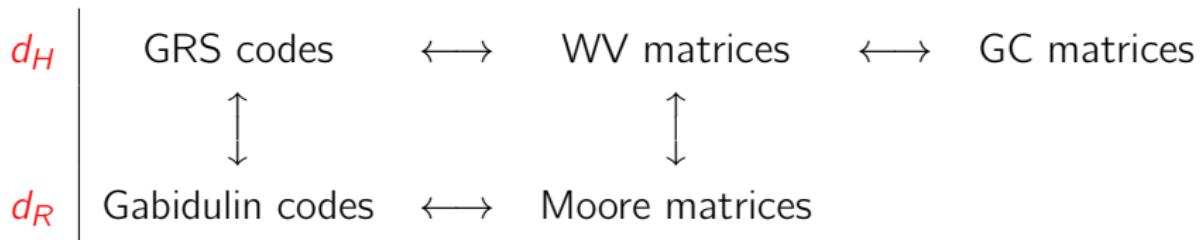
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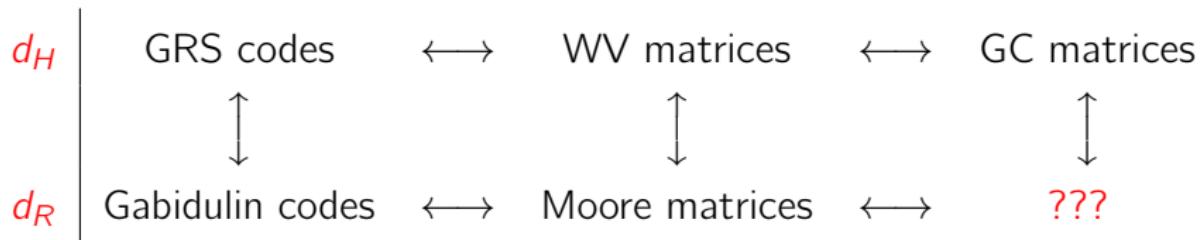
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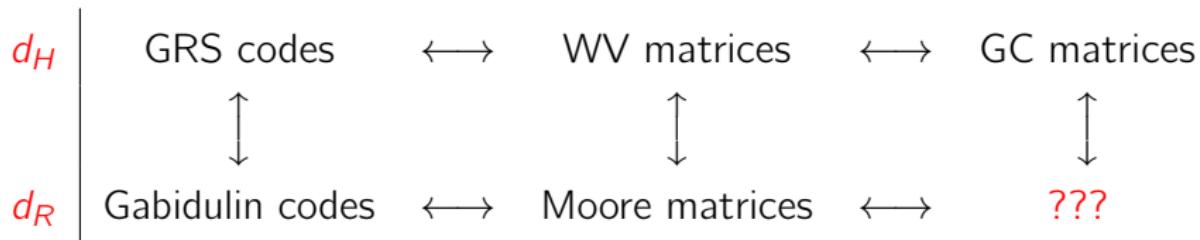
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Hamming vs Rank Distance: Recap



(Possible) definition of a q -analogue of Cauchy matrices!

GRC Matrix

Let $\gamma \in \mathbb{F}_{q^m}$ such that $\text{Tr}(\gamma) \neq 0$ and s be an integer coprime to m . We define the map

$$\begin{aligned}\pi : \mathbb{F}_{q^m} &\longrightarrow \mathbb{F}_{q^m} \\ \alpha &\longmapsto -\frac{1}{\text{Tr}(\gamma)} \sum_{i=0}^{m-2} \left(\sigma^{i+1}(\gamma) \sum_{j=0}^i (\sigma^j(\alpha)) \right)\end{aligned}$$

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- (1) $\alpha_1, \dots, \alpha_k \in \mathbb{F}_{q^m}$, \mathbb{F}_q -linearly independent,
- (2) $\beta_1, \dots, \beta_{n-k} \in \mathbb{F}_{q^m}$, \mathbb{F}_q -linearly independent,
- (3) $\beta_1, \dots, \beta_{n-k} \in \langle \alpha_1, \dots, \alpha_k \rangle_{\mathbb{F}_q}^\times$,
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The matrix $X \in \mathbb{F}_{q^m}^{k \times (n-k)}$ defined by

$$X_{i,j} = \pi(\alpha_i \beta_j) + c_{i,j}$$

is a **Generalized Rank-Cauchy (GRC) Matrix**.

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Standard Form of Gabidulin Codes

Theorem 1 [N. '18]

There is a 1-1 correspondence between GRC matrices and Gabidulin codes given by

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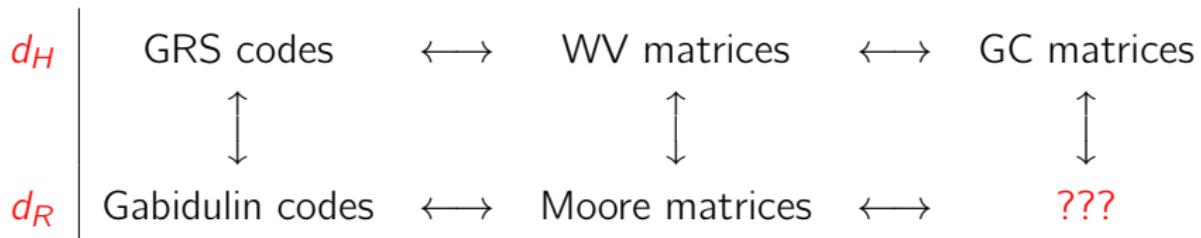
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For every $A \in \text{GL}_k(\mathbb{F}_q)$, $B \in \text{GL}_{n-k}(\mathbb{F}_q)$,
every minor of the matrix AXB is non-zero! (*)

Standard form of Gabidulin codes

Standard form of Gabidulin codes



Standard form of Gabidulin codes

